

Model Answer + Grading Rubric

Answer to exercise 01 : (COMPARISON OF QUEUES : 05 Marks)

We have $\rho_1 = \rho_2 = \rho = \frac{\lambda}{\mu}$.

1. Probability of an Empty System :

$$P_{0M/M/1} = 1 - \rho$$

0.5

$$P_{0M/M/2} = \left[\frac{(2\rho)^2}{2(1-\rho)} + 1 + 2\rho \right]^{-1} = \left[\frac{2\rho^2}{1-\rho} + \frac{1-\rho}{1-\rho} + \frac{2\rho(1-\rho)}{1-\rho} \right]^{-1} = \left[\frac{2\rho^2 + 1 - \rho + 2\rho(1-\rho)}{1-\rho} \right]^{-1} = \left[\frac{1+\rho}{1-\rho} \right]^{-1}$$

$$P_{0M/M/2} = \frac{1-\rho}{1+\rho}$$

0.5

2. We notice that $P_{0M/M/2} = \frac{P_{0M/M/1}}{1+\rho}$. $\rho > 0 \Rightarrow 1 + \rho > 1$. Hence,

$$P_{0M/M/2} < P_{0M/M/1}$$

0.5

3. Average Number of Waiting Customers in the System :

$$\bar{Q}_{M/M/1} = \frac{\rho^2}{1-\rho}$$

0.5

$$\zeta = \frac{2\rho^2}{1-\rho} \frac{1-\rho}{1+\rho} = \frac{2\rho^2}{1+\rho}. \text{ So, } \bar{Q}_{M/M/2} = \frac{\zeta\rho}{1-\rho} = \frac{2\rho^3}{(1-\rho)(1+\rho)}. \text{ Hence, } \bar{Q}_{M/M/2} = \frac{2\rho^3}{1-\rho^2}$$

$$\bar{Q}_{M/M/2} = \frac{2\rho^3}{1-\rho^2}$$

1.0

0.5

4. We notice that $\bar{Q}_{M/M/2} = \frac{\rho}{1+\rho} \bar{Q}_{M/M/1}$. Since $\rho < 1 + \rho \Rightarrow \frac{\rho}{1+\rho} < 1$. We conclude that

$$\bar{Q}_{M/M/2} < \bar{Q}_{M/M/1}$$

5. Average Number of Customers in the System :

$$\bar{N}_{M/M/1} = \frac{\rho}{1-\rho}$$

0.5

$$\bar{N}_{M/M/2} = \bar{Q}_{M/M/2} + \bar{R}_{M/M/2} = \frac{2\rho^3}{1-\rho^2} + 2\rho = \frac{2\rho^3 + 2\rho(1-\rho^2)}{1-\rho^2}. \text{ Hence, } \bar{N}_{M/M/2} = \frac{2\rho}{1-\rho^2}$$

$$\bar{N}_{M/M/2} = \frac{2\rho}{1-\rho^2}$$

0.5

6. We notice that $\bar{N}_{M/M/2} = \frac{2}{1+\rho} \bar{N}_{M/M/1}$. We have $\rho < 1 \Rightarrow 1 + \rho < 2 \Rightarrow \frac{2}{1+\rho} > 1$.

$$\bar{N}_{M/M/2} > \bar{N}_{M/M/1}$$

0.5

Answer to exercise 02 : (QUEUEING SYSTEM : 05 Marks)

1. We have $\lambda = 5$ requests/s and $\mu = 6$ requests/s. The appropriate model for the system is $M/M/1/N$ because :

- The arrival process is Poisson (statement).
- Exponential service time (statement).
- The number of servers is $m = 1$ (statement).
- Service discipline is *FIFO* (default).
- Finite capacity (N requests in the system/limited memory size).
- Infinite population size (default).

1.0

1.0

2. The minimal memory space required to reach the objective, if each request needs 1500 Bytes is :

$$\rho = \frac{\lambda}{\mu} = \frac{5}{6} \approx 0,83333333 \neq 1, P_0 = \frac{1-\rho}{1-\rho^{N+1}}, P_N = \frac{(1-\rho)\rho^N}{1-\rho^{N+1}}$$

The effective request rate accepted by the server is λ_e . To satisfy at least 90% of the received requests, $\frac{\lambda_e}{\lambda} \geq 0.9$ must hold.

$$\frac{\lambda_e}{\lambda} \geq 0.9 \Rightarrow \frac{\lambda(1-P_N)}{\lambda} \geq 0.9 \Rightarrow 1 - P_N \geq 0.9 \Rightarrow 1 - \frac{(1-\rho)\rho^N}{1-\rho^{N+1}} \geq 0.9 \Rightarrow \frac{1-\rho^N}{1-\rho^{N+1}} \geq 0.9$$

$$\text{Hence : } \frac{\lambda_e}{\lambda} \geq 0.9 \Rightarrow (1 - 0.9\rho)\rho^N \leq 0.1 \Rightarrow \rho^N \leq \frac{0.1}{1-0.9\rho} \Rightarrow \rho^N \leq 0,4 \Rightarrow N \ln(\rho) \leq \ln(0,4)$$

$$\frac{\lambda_e}{\lambda} \geq 0.9 \Rightarrow N \ln\left(\frac{5}{6}\right) \leq \ln(0,4) \Rightarrow N \geq \frac{\ln(0,4)}{\ln\left(\frac{5}{6}\right)} \quad (\ln(\rho) < 0 \text{ because } \rho < 1)$$

1.0

Finally, $N \geq 5,025685101562$ and at least $N = 6$. Needed memory = $N \times 1500 = 9600 \text{ Bytes}$.

3. The average number of daily processed requests :

Since there is a reject and only part of the flow enters the system, we use the effective occupation rate here :

$$P_N = P_6 = \frac{(1-\rho)\rho^N}{1-\rho^{N+1}} = 0.07742392,$$

$$\rho_e = \frac{\lambda_e}{\mu} = \frac{\lambda(1-P_N)}{\mu} = \frac{5 \times (1-0.07742392)}{6} = 0.7688133947$$

Duration of activity per day (24 hours) : $24 \times \rho_e = 24 \times 0.7688133947 = 18,4515214730$ hours = 1107,0912883839 minutes = 66425,4773030345 seconds.

$$Nb = \left\lfloor \frac{\text{Activity duration}}{S} \right\rfloor = \left\lfloor \text{Activity duration} \times \mu \right\rfloor = \left\lfloor 66425,4773030345 \times 6 \right\rfloor = \left\lfloor 398552,8638182071 \right\rfloor \Rightarrow$$

$$Nb = 398552 \text{ requests} \quad \leftarrow 1.0$$

4. The average number of waiting requests :

$$P_0 = \frac{1-\rho}{1-\rho^{N+1}} = 0.231186605289$$

$$\bar{Q} = \frac{(m\rho)^m \rho P_0}{m!(1-\rho)^2} [1 - \rho^{N-m} - (N-m)(1-\rho)\rho^{N-m}]$$

$$\bar{Q} = \frac{0.83333333^2 \times 0.231186605289}{1!(1-0.83333333)^2} [1 - 0.83333333^5 - 5(1-0.83333333)0.83333333^5]$$

$$\text{Hence : } \bar{Q} = 1.521349183146 \text{ requests} \quad \leftarrow 1.0$$

Answer to exercise 03 : (QUEUEING NETWORK : 10 Marks)

Nous avons : $\gamma = 5, m_1 = +\infty, m_2 = 2, m_3 = 1, \mu_1 = 1, \mu_2 = 5, \mu_3 = 10$.

1. Internal routing probability matrix :

$$\begin{pmatrix} \frac{1}{3} & \frac{1}{3} & \frac{1}{3} \\ \alpha & 0 & 0 \\ \alpha & 0 & 0 \end{pmatrix}$$

External routing probability matrix :

$$\begin{pmatrix} 0 \\ 1-\alpha \\ 1-\alpha \end{pmatrix}$$

$\leftarrow 1.0$

$\leftarrow 1.0$

2. Effective arrival rates λ_i :

$$\begin{cases} \lambda_1 = \gamma + \frac{1}{3}\lambda_1 + \alpha\lambda_2 + \alpha\lambda_3 \\ \lambda_2 = \frac{1}{3}\lambda_1 \\ \lambda_3 = \frac{1}{3}\lambda_1 \end{cases} \Rightarrow \begin{cases} \frac{2}{3}\lambda_1 = 4 + \frac{2\alpha}{3}\lambda_1 \\ \lambda_2 = \lambda_3 = \frac{1}{3}\lambda_1 \end{cases} \Rightarrow \begin{cases} \frac{2(1-\alpha)}{3}\lambda_1 = 4 \\ \lambda_2 = \lambda_3 = \frac{1}{3}\lambda_1 \end{cases}$$

$$\Rightarrow \begin{cases} \lambda_1 = \frac{6}{1-\alpha} \\ \lambda_2 = \lambda_3 = \frac{2}{1-\alpha} \end{cases} \Rightarrow \begin{cases} \lambda_1 = \frac{6}{1-\alpha} \\ \lambda_2 = \frac{2}{1-\alpha} \\ \lambda_3 = \frac{2}{1-\alpha} \end{cases} \quad \leftarrow 1.5$$

Note that $\alpha = 1$ cannot be a solution.

3. Values of α which ensure the stability of the network :

$$\begin{cases} FA_1 \text{ Stable} \\ \wedge \\ FA_2 \text{ Stable} \\ \wedge \\ FA_3 \text{ Stable} \end{cases} \Rightarrow \begin{cases} \rho_1 = 0 < 1 \\ \wedge \\ \rho_2 = \frac{1}{4(1-\alpha)} < 1 \\ \wedge \\ \rho_3 = \frac{1}{6(1-\alpha)} < 1 \end{cases} \Rightarrow \begin{cases} 1-\alpha > \frac{1}{4} \\ \wedge \\ 1-\alpha > \frac{1}{6} \end{cases} \Rightarrow \begin{cases} \alpha < \frac{3}{4} \\ \wedge \\ \alpha < \frac{5}{6} \end{cases}$$

$$\text{network stable} \Rightarrow \begin{cases} 0 \leq \alpha < \frac{3}{4} \\ \wedge \\ 0 \leq \alpha < \frac{5}{6} \end{cases}$$

We conclude that for the network to be stable, α must verify the following condition :

$$0 \leq \alpha < \frac{3}{4} \quad \text{or} \quad \alpha \in [0, \frac{3}{4}[\quad \leftarrow 1.0$$

For $\alpha = \frac{1}{2}$:

$$\begin{cases} \lambda_1 = 12 \\ \lambda_2 = 4 \\ \lambda_3 = 4 \end{cases} \Rightarrow \begin{cases} \rho_1 = 0 \\ \rho_2 = \frac{\lambda_2}{m_2\mu_2} = \frac{4}{8} \\ \rho_3 = \frac{\lambda_3}{m_3\mu_3} = \frac{4}{2 \times 6} \end{cases} \Rightarrow \begin{cases} \rho_1 = 0 \\ \rho_2 = \frac{1}{2} = 0.5 \\ \rho_3 = \frac{1}{3} = 0.3333333333 \end{cases}$$

Queue 1 : of type $M/M/+ \infty$. **Queue 2 :** de type $M/M/2$. **Queue 3 :** de type $M/M/1$.

4. The average number of customers waiting in each queue and in the network :

(a) $\overline{Q}_1 = 0$ 0.25

(b) $\overline{Q}_2 = \frac{\rho_2^2}{1-\rho_2} = \frac{0.5^2}{1-0.5} \Rightarrow \overline{Q}_2 = \frac{1}{2} = 0.5$ 0.25

(c) $P_0(3) = \left[\frac{(m_3 \rho_3)^{m_3}}{m_3!(1-\rho_3)} + \sum_{k=0}^{m_3-1} \frac{(m_3 \rho_3)^k}{k!} \right]^{-1} = \left[\frac{(2 \times \frac{1}{3})^2}{2!(1-\frac{1}{3})} + 1 + 2 \times \frac{1}{3} \right]^{-1} = \left[\frac{1}{3} + 1 + 2 \times \frac{1}{3} \right]^{-1}$
 $P_0(3) = \left[\frac{1}{3} + \frac{2}{3} + \frac{2}{3} \right]^{-1} \Rightarrow P_0(3) = \frac{1}{2} = 0.5$ 0.25

$\zeta_3 = \frac{(m_3 \rho_3)^{m_3}}{m_3!(1-\rho_3)} P_0(3) = \frac{(2 \times \frac{1}{3})^2}{2!(1-\frac{1}{3})} P_0(3) = \frac{1}{3} \times \frac{1}{2} \Rightarrow \zeta_3 = \frac{1}{6} = 0.1666666666$ 0.25

$\overline{Q}_3 = \frac{\zeta_3 \rho_3}{1-\rho_3} = \frac{1 \times 3}{6 \times 3 \times 2} \Rightarrow \overline{Q}_3 = \frac{1}{12} = 0.0833333333$ 0.25

(d) $\overline{Q}_R = \sum_{i=1}^3 \overline{Q}_i = 0 + \frac{1}{2} + \frac{1}{12} \Rightarrow \overline{Q}_R = \frac{7}{12} = 0.5833333333$ 0.25

5. The average number of customers in each queue and in the network :

(a) $\overline{N}_1 = \overline{R}_1 = \frac{\lambda_1}{\mu_1} \Rightarrow \overline{N}_1 = \frac{12}{5} = 2.4$ 0.25

(b) $\overline{N}_2 = \overline{Q}_2 + \overline{R}_2 = 0.5 + 0.5 \Rightarrow \overline{N}_2 = 1$ 0.25

(c) $\overline{N}_3 = \overline{Q}_3 + \overline{R}_3 = \frac{1}{12} + \frac{2}{3} \Rightarrow \overline{N}_3 = \frac{3}{4} = 0.75$ 0.25

(d) $\overline{N}_R = \sum_{i=1}^3 \overline{N}_i = \frac{12}{5} + 1 + \frac{3}{4} = \frac{48+20+15}{20} \Rightarrow \overline{N}_R = \frac{83}{20} = 4.15$ 0.25

6. The average residence time in each queue and in the network :

(a) $\overline{T}_1 = \frac{1}{\mu_1} = \frac{1}{5} \Rightarrow \overline{T}_1 = 0.2$ 0.25

(b) $\overline{T}_2 = \frac{\overline{N}_2}{\lambda_2} = \frac{1}{4} \Rightarrow \overline{T}_2 = 0.25$ 0.25

(c) $\overline{T}_3 = \frac{\overline{N}_3}{\lambda_3} = \frac{3}{4 \times 4} \Rightarrow \overline{T}_3 = \frac{3}{16} = 0.1875$ 0.25

(d) $\overline{T}_R = \frac{\overline{N}_R}{\lambda_R} = \frac{\overline{N}_R}{\sum_{i=1}^3 \gamma_i} = \frac{83}{20 \times 4} \Rightarrow \overline{T}_R = \frac{83}{80} = 1.0375$ 0.25

7. The average waiting time in each queue and in the network :

(a) $\overline{W}_1 = 0$ 0.25

(b) $\overline{W}_2 = \frac{\overline{Q}_2}{\lambda_2} = \frac{1}{2 \times 4} \Rightarrow \overline{W}_2 = \frac{1}{8} = 0.125$ 0.25

(c) $\overline{W}_3 = \frac{\overline{Q}_3}{\lambda_3} = \frac{1}{12 \times 4} \Rightarrow \overline{W}_3 = \frac{1}{48} = 0.0208333333$ 0.25

(d) $\overline{W}_R = \frac{\overline{Q}_R}{\lambda_R} = \frac{\overline{Q}_R}{\sum_{i=1}^3 \gamma_i} = \frac{0 + \frac{1}{2} + \frac{1}{12}}{4} \Rightarrow \overline{W}_R = \frac{7}{48} = 0.1458333333$ 0.25

8. The probability that the network is not empty :

$Pr(\text{Network not empty}) = 1 - Pr(\text{Network is empty}) = 1 - (Pr(\text{Queue}_1 \text{ empty and Queue}_2 \text{ empty and Queue}_3 \text{ empty}))$
 $= (1 - P_0(\text{Queue}_1) \times P_0(\text{Queue}_2) \times P_0(\text{Queue}_3))$

$Pr(\text{Network not empty}) = 1 - e^{-\frac{\lambda_1}{\mu_1}} \times P_0(2) \times P_0(3) = 1 - e^{-\frac{12}{5}} \times (1 - \rho_2) \times P_0(3)$
 $= 1 - 0.0907179532 \times 0.5 \times 0.5 = 1 - 0.0226794883$

Hence : $Pr(\text{Network not empty}) = 0.9773205116$ 1.00